APNIC 56 Routing Security SIG

Autonomous System Relationship Authorization

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- 1. Background
- 2. Autonomous System Provider Authorization (ASPA)
- 3. Autonomous System Relationship Authorization (AS<u>R</u>A)

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Background









RFC7908 : Problem Definition and Classification of BGP Route Leaks

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- ASPA (Autonomous System Provider Authorization): Detect route leakage and part of origin-forged attack based on business relationship
- **RPKI ASPA record:** ASPA (ASx, [ASy1, ASy2, ...]), where Asx is the registration AS and ASy is the provider AS. Store data in RPKI repo.



• **Hop-check:** Given an AS pair (ASx, ASy), Hop-check() function checks whether ASy is the provider of ASx based the ASPA data and returns Provider+, Not provider+, or No attestation.



- Upstream check: verify the route received from customer, lateral peer, RS, or RS-client. Simply,
 - > Valid: Each pair should be Provider+
 - > Invalid: Not provider+ exists
 - > Unknown: Others



Example 1: AS5 verifies path [AS1, AS2, AS3, AS4] and gets valid/unknown result



Example 2: AS5 verifies path [AS1, AS2, AS3, AS4] and gets invalid result



- Downstream check: verify the route received from transit or mutual transit provider. Simply,
 - > Valid: The first valid up ramp and the last valid down ramp composite the complete path
 - > Invalid: More than two invalid hops between the first valid up ramp and the last valid down ramp

> Unknown: Others



Example 3: AS8 verifies the path and gets valid/unknown result

The first valid up ramp: [1, 2, 3, 4] The last valid down ramp: [5, 6, 7]



Example 4: AS8 verifies the path and gets invalid result

The first valid up ramp: [1, 2, 3, 4] The last valid down ramp: [7]

ASPA Analysis

- ASPA can effectively detect invalid paths resulted by route leakage and part of originforged attacks, if the ASPA data are correctly and adequately registered
- ASPA is a promising mechanism

- There are also some remaining challenges of ASPA (cover more scenarios?):
 - > How to guarantee **the correctness of ASPA data registration**
 - > How to improve the ability of **detecting path hijacks (fake links)**
 - > How to improve the validation benefits **under partial registration**
 - > How to cope with complex scenarios such as legitimate valley-paths, hybrid/partial transit relationships

Case 1: Cannot cross-check the correctness of registration



- AS1 registers: AS2, AS3, and AS4 are my providers
- But, AS2 is not its provider
- The incorrectness cannot be detected by the ASPA data registered by AS2

Case 2: Cannot detect hijack introduced by provider

 Route received by AS108: {Origin ASN=106, AS path = [107, 106]}



• ASPA result: valid

• Analysis: In the current ASPA design, any downstream paths with length no longer than 2 will be considered valid

<pre>root@ip-s8:~/gengnan/bgp/open flags: * = Valid, > = Selecto S = Stale, E = Error origin validation state: N = aspa validation state: ? = un origin: i = IGP, e = EGP, ? =</pre>	ed, I = via IBGP, not-found, V = va nknown, V = valid,	Ă = Annou lid, ! =	unced, invalid
<pre>flags vs destination *> V-V 192.106.0.0/24 *> V-V 192.107.0.0/24 AI*> V-? 192.108.0.0/24</pre>	10.107.10.154	100 100	med aspath origin 0 107 106 i 0 107 i 0 i

• Expected result: invalid-hijack

ASPA records:

	AS106	AS107	
ASPA	provider-set = {105}	provider-set = {0}	

Case 3: Cannot detect leakage with partial registration

 Route received by AS108: {Origin ASN=106, AS path = [107, 106]}



ASPA records:

	AS106	AS107
ASPA	No data	provider-set = {108}

ASPA result: unknown

• Analysis: When do upstream check, hop-check(106, 107) returns No Attestation because AS106 has no record

flags: * = Val S = Sta origin validat aspa validatio	<pre>jengnan/bgp/openbg id, > = Selected, ile, E = Error ion state: N = no on state: ? = unkn iP, e = EGP, ? = I</pre>	I = via IBGP, A t-found, V = val own, V = valid,	.id, ! = :	nced, invalid
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*> V-? 192.	106.0.0/24	10.107.10.154	100	0 107 106 i
*> V-V 192.	107.0.0/24	10.107.10.155	100	0 107 i
AI*> V-? 192.	108.0.0/24	0.0.0.0	100	0 i

• Expected result: invalid

Case 4: Cannot cope with complex scenarios

 Route received by AS108: {Origin ASN=106, AS path = [107, 106]}



ASPA records:

	AS106	AS107
ASPA	provider-set = {105}	

ASPA result: invalid

• Analysis: In the current ASPA design, legitimate valley-paths, hybrid/partial transit relationships cannot be coped with

<pre>root@ip-s8:~/gengnan/bgp/ope flags: * = Valid, > = Select S = Stale, E = Error origin validation state: N = aspa validation state: ? = u origin: i = IGP, e = EGP, ?</pre>	ed, I = via IBGP, not-found, V = va nknown, V = valid,	Ă = Anno lid, ! =	unced, invalid
flags vs destination	gateway	lpref	med aspath origin
*> V-! 192.106.0.0/24	10.107.10.154	100	0 107 106 i
*> V-V 192.107.0.0/24	10.107.10.155	100	0 107 i
AI*> V-? 192.108.0.0/24	0.0.0.0	100	0 i

• Expected result: valid

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Autonomous System <u>Relationship</u> Authorization (AS<u>R</u>A)

• Preliminary idea: Register more relationship information

Level	Feature	Content	Description
1	Mandatory	Providers	Same as ASPA
2	Optional	Neighbors	All AS adjacencies, similar to [1]
3	Optional	All normal and complex relationships	Customer, lateral peer, hybrid, partial, valley-path, etc.

• [1] draft-huston-sidr-aao-profile-03. A Profile for AS Adjacency Attestation Objects. Geoff Huston , George G. Michaelson

Illustration of AS<u>R</u>A Record



Data type	Requirement	Record data
provider set	Mandatory	[1, 2, 3]
other neighbor set	Optional	[4, 5, 6, 7]
customer set		[3, 5]
lateral peer set		[4]
partial transit set	Optional	[6]
hybrid set		[(7,tag1:P, tag2:C)], tag means geographic location
valley-path set		[[1, 2]], means 1, X, and 2 can form a valley-path

What Can be Done using ASRA Data

• Usage 1: Cross-check the correctness of registration

> X registers Y as the provider, but Y registers X as the peer. A conflict occurs.

> X states Y is the provider, but X does not appear in the neighbor set of Y. A conflict occurs.

• Usage 2: Identify fake links

> An AS can register all neighboring ASes. Fake links in the AS path can be detected and considered as hijacking.

• Usage 3: Hop-check(Y, X) for Hop-check(X, Y) under partial registration

> In Hop-check(X, Y), if X does not register data, you can use the ASRA data registered by Y to verify the Xto-Y business relationship.

• Usage 4: Cope with complex scenarios

> Legitimate valley-path, hybrid/partial relationship

ASRA Implementation based on OpenBgpd

- Modify the open source project OpenBgpd for implementing the ASRA prototype
- The config file is extended. Three new kinds of records are supported:
 - > Neighbor set
 - > Customer set

```
> Peer set
```

aspa-set {							
customer-as	107	provider-as	{106,	108}	customer-set	<pre>{0} peer-set</pre>	{0}
customer-as	106	provider-as	{105}	neigl	hbor-set {111,	, 222}	
}				-			

- The verification is modified, and the ASRA records can be fully used.
- The complex relationships and legitimate valley-path are not supported. RTR protocol is not extended.
- Experiment: 3 servers, Ubuntu 20.04.5 LTS, named as AS106, AS107, and AS108, respectively



Case 1: Can detect hijack introduced by provider

 Route received by AS108: {Origin ASN=106, AS path = [107, 106]}



ASPA & ASRA records:

	AS106	AS107
ASPA	provider-set = {105}	
ASRA	provider-set = {105} neighbor-set = {111, 222}	
ASRA	provider-set = {105} customer-set = {111} peer-set = {222}	

ASPA result: valid

• Analysis: In the current ASPA design, any downstream paths with length no longer than 2 will be considered valid

• NEW result: invalid-hijack

• Analysis: Since all neighbor links are known, the fake link can be easily detected

flags: origin aspa N	<pre>ip-s8:~/gengnan/bgp/openb : * = Valid, > = Selected S = Stale, E = Error n validation state: N = n validation state: ? = unk n: i = IGP, e = EGP, ? =</pre>	, I = via IBGP, ot-found, V = va nown, V = valid,	A = Annou Lid, ! =	nced,
	vs destination V-!! 192.106.0.0/24	gateway 10.107.10.154		med aspath origin 0 107 106 i
*>	V-V 192.107.0.0/24 V-? 192.108.0.0/24	10.107.10.155 0.0.0.0		

Case 2: Can detect leakage with partial registration

 Route received by AS108: {Origin ASN=106, AS path = [107, 106]}



ASPA & ASRA records:

	AS106	AS107
ASPA	无	provider-set = {108}
ASRA	无	provider-set = {108} customer-set = {0} peer-set={0}

ASPA result: unknown

• Analysis: When do upstream check, hop-check(106, 107) returns No Attestation because AS106 has no record

• NEW result: invalid

 Analysis: Backward verification is used. When AS106 does not have registration data, AS107 can be used to register customer and peer data for verification. Only provider data (sibling cannot be identified) or provider and neighbor data cannot be used for backward verification.

<pre>root@ip-s8:~/gengnan/bgp/openbgpd-portable # bgpctl show rib flags: * = Valid, > = Selected, I = via IBGP, A = Announced, S = Stale, E = Error origin validation state: N = not-found, V = valid, ! = invalid aspa validation state: ? = unknown, V = valid, ! = invalid, !! = invalid-hijack origin: i = IGP, e = EGP, ? = Incomplete</pre>					
flags vs destination	gateway	lpref	med aspath origin		
*> V-!! 192.106.0.0/	24 10.107.10.154	100	0 107 106 i		
*> V-V 192.107.0.0/2	4 10.107.10.155	100	0 107 i		
AI*> V-? 192.108.0.0/2	4 0.0.0.0	100	0 i		

Case 3: Can detect hijack with partial registration

 Route received by AS108: {Origin ASN=106, AS path = [107, 106]}



ASPA & ASRA records:

	AS106	AS107
ASPA	无	provider-set = {108}
ASRA	无	provider-set = {108} neighbor-set = {0}
ASRA	无	provider-set = {108} customer-set = {0} peer-set={0}

ASPA result: unknown

• Analysis: When do upstream check, hop-check(106, 107) returns No Attestation because AS106 has no record

• NEW result: invalid-hijack

192.108.0.0/24

 Analysis: Backward verification is used. When AS 106 has no registration data, the data registered by AS 107 can be used for verification to identify hijacking

<pre>root@ip-s8:~/gengnan/bgp/openbgpd-portable # bgpctl show rib flags: * = Valid, > = Selected, I = via IBGP, A = Announced, S = Stale, E = Error origin validation state: N = not-found, V = valid, ! = invalid aspa validation state: ? = unknown, V = valid, ! = invalid, !! = invalid-hijack origin: i = IGP, e = EGP, ? = Incomplete</pre>					
flags vs destination *> V-!! 192.106.0.0/24 *> V-V 192.107.0.0/24		lpref 100 100			

100

0 i

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Possible Extensions for ASRA

- Possible extensions for ASRA deployment
 > ASRA Profile
 Registration is challenging
 - > SLURM for local data provision
 - > RTR protocol for ASRA record synchronization
 - > ASRA verification (internal implementation)



Deployment Consideration

- Globally public data registration is challenging. Some ASes may be concerned about **privacy**.
 - > There are some analysis on why privacy problem is not much important in many cases [1]
 - > Selective registration. ASes can choose to register neighbor set, detailed relationships, etc. Neighbor set induces relatively less privacy concerns.
 - > Regional registration. ASRA data can be registered and used with a region, which can efficiently prevent Type1 and Type2 origin-forged attacks.
- ASRA can improve deployment benefits under partial registration.
 - > Generally, ASPA has less deployment benefits than ROA under the same registration ratio
- [1] Cohen, Avichai, et al. "Jumpstarting BGP security with path-end validation." Proceedings of the 2016 ACM SIGCOMM Conference. 2016.

Conclusion

- ASRA: AS can optionally register more detailed AS relationships, and ASPA can be enhanced.
 > Usage 1: Cross-check the correctness of registration
 > Usage 2: Identify fake links
 > Usage 3: Improve deployment benefits under partial registration
 - > Usage 4: Cope with complex scenarios

• Limitation: Can detect route leakage and fake link hijacking, but cannot prevent and identify AS path tampering (BGPsec for path protection)

Thank you!