IPv6 (Internet Protocol version 6)

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=> What's IPv6

Address Architecture Plug & play Domain Name System Transition The current status of IPv6

Problems of IPv4

Exhaustion of IPv4 addresses

○4 bytes = 4.3 billion

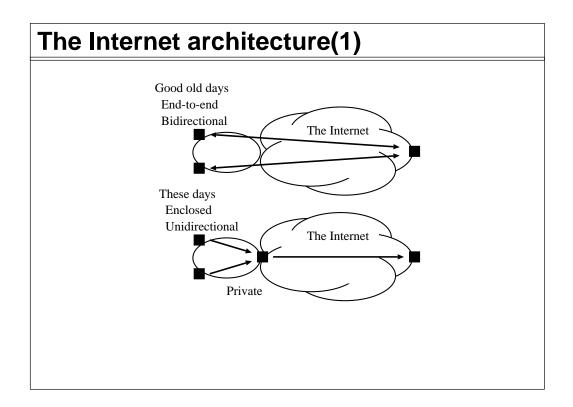
- ▷ Much less than the human population (6.1 billion)
- ° Will be exhausted in around 2008
- °Registries are allocating IPv4 addresses by severe policy
 - ▷Nobody can obtain enough IPv4 addresses

□ Increment of routing information

- $^{\rm O} {\rm Routing}$ information cannot be aggregated effectively
 - Unaggregatable address assignments
- °80,000 entries at present
- **OBurden for backbone routers**
 - Unstability, accidents

□ Proliferation of NAT

- ^OBreaking the Internet architecture
- Enclosure of users



The Internet architecture(2) □ Flat Internet • Bidirectional communication • End-to-end communication • True communication infrastructure • Much easier to deploy inventive new applications □ Patched Internet by NAT • Unidirectional communication • Enclosed communication • Single point of failure • Evolution of application are suppressed by NAT • Accounting from servers is impossible

IPv6

□Address extension

○16 bytes = 3.4 x 10^38

- •65,536 subnets (minimum) for everybody (/48)
 - \triangleright Class A (in the old term) per site

□Starting over

- •Some technologies are mandatory
 - ▶ Plug and play
 - ▷ End-to-end security (i.e. IPsec)
- ○Aggregatable global addresses from the start ▷Reducing external routing information to 8,192

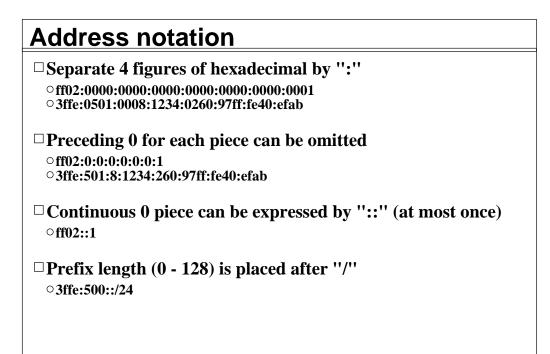
□ Paradigm chage for applications

- °End-to-end and bidirectional communication
- IPv6 is a NAT-free world
 - $^{\triangleright}$ Cellular phone, automobiles, home network, game machines, ...

What's IPv6

=> Address Architecture

Plug & play Domain Name System Transition The current status of IPv6



Address block

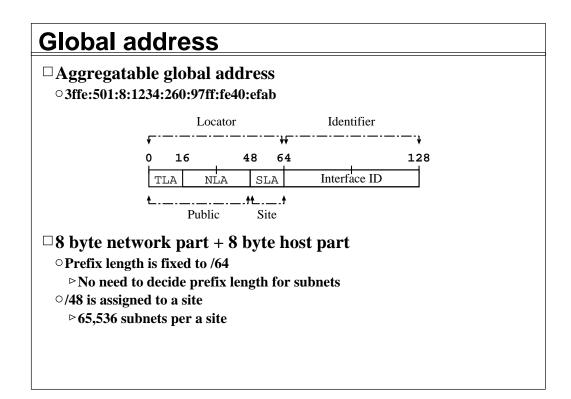
□ The address space of IPv6 is divided into 8 blocks (/3)

 $^{\circ}3$ bits is not friendly to hexadecimal

 $^{\circ}$ Let's consider /4 (i.e. 16 blocks)

□ The first hexadecimal figure:

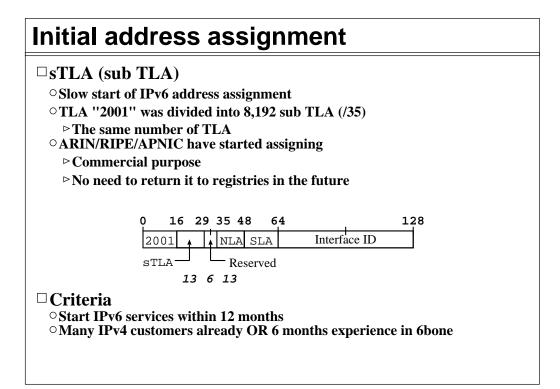
- 0,1 Special (e.g loopback)
- 02,3 Global address (aggregatable global address)
- **04,5** Not assigned
- **°6,7** Not assigned
- **08,9** Not assigned
- oa, b Not assigned
- oc, d Not assigned
- $^{\rm O}$ e, f $\,$ Link-local, site-local, multicast

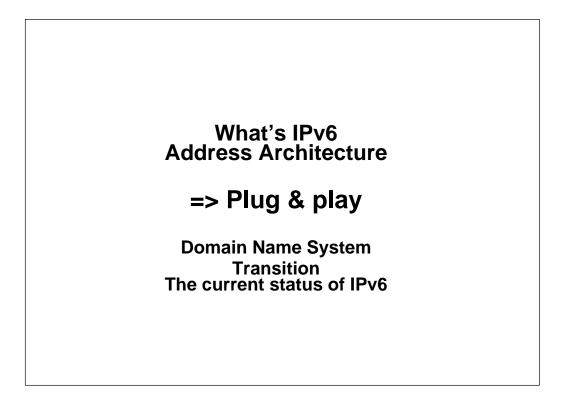


Aggregator

○65,536 subnets per site

□ TLA (Top Level Aggregator)
0/16
○ Big ISPs or IXes
○ 8,192 (16 - 3 = 13 bit)
○ MUST announce /16 routing information to other TLAs
□ NLA (Next Level Aggregator)
0/17 - /48
○ Medium or small ISPs
▷ NLA1, NLA2,...
○ Final NLA is a site
□ SLA (Site Level Aggregator)
0/49 - /64
○ Subnet number in a site





Two stories

□ Dilemma of dentists

 $^{\rm O}$ Well-educated but a beginner of network

 $^{\rm O}$ They believe that computers can start communicating just after they open

boxes

□Nightmare of network managers

- °100 computers arrived on Friday
- $^{\rm O}\mbox{They}$ have to get computers ready for the Internet by Monday

Address auto-configuration Lower 8 bytes from a MAC address Link-local address Ethernet address Upper 8 bytes from a router Global address DHCP server is not necessary Stateless address allocation

Auto-configuration of link-local

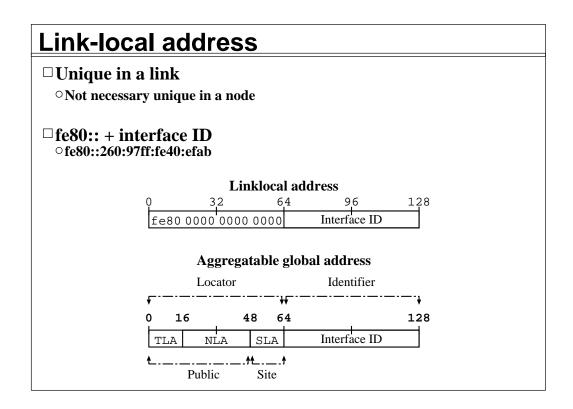
□ MAC address (e.g Ethernet) • IEEE 802 address (6 bytes) • 00:60:97:40:ef:ab

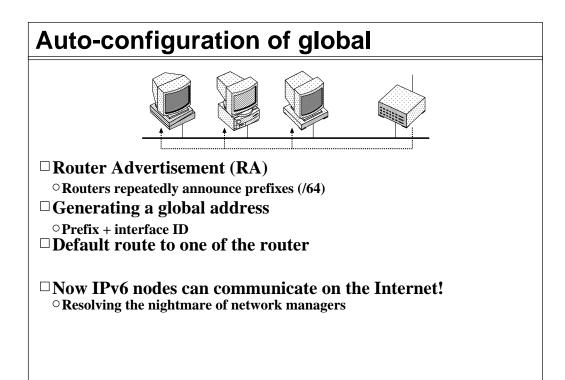
□ Converting MAC address to interface ID

○ EUI 64 address (8 bytes)
 ○ 260:97ff:fe40:efab
 ▷ 00:60:97 + ff:fe + 40:ef:ab
 ▷ Negate one bit

□ Generating a link-local address

□ Now IPv6 nodes can communicate within the link! ○ Resolving the dilemma of dentists





Renumbering

□ **Precondition**

○All IPv6 nodes have address auto-configuration

 $^{\odot}\textbf{Each}$ IPv6 node can obtain multiple IPv6 addresses

°IPv6 addresses have two timers

□ Switching ISP A to ISP B

•A site is connected to ISP A

▷An old address

 $^{\rm O}$ Connect to ISP B, then a new prefix is announced

▷ The old address, a new address

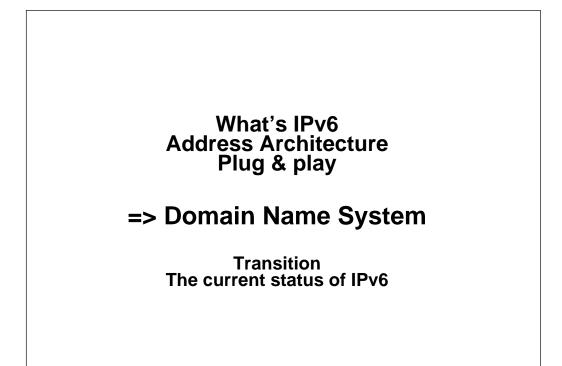
 $^{\circ}$ The first timer is expired

▶ The old address is not used for further communication

 $^{\rm O}$ The second timer is expired

▷ The new address

 $^{\rm O}\textsc{Disconnect}$ the leased to ISP A



Using IPv6 applications

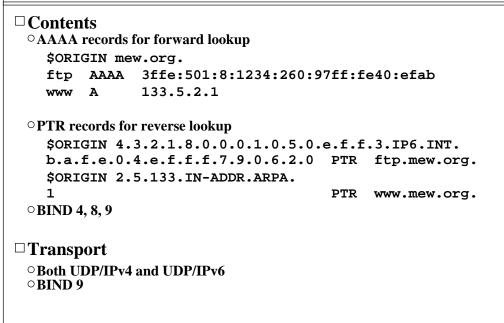
□ Typical users specify "host name" to applications °(e.g) ftp ftp.mew.org

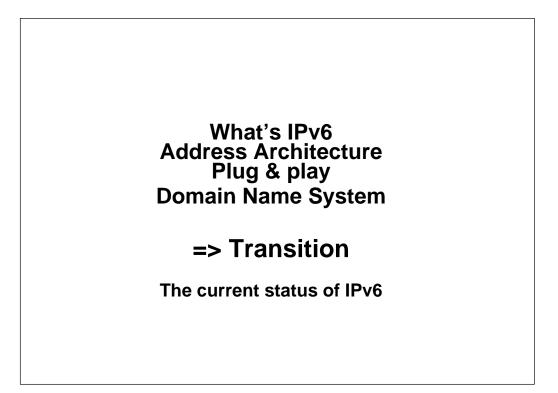
□ Users are not aware which they are using, IPv4 or IPv6 °(e.g) "ftp ftp.mew.org" may use IPv4 °(e.g) "ftp ftp.mew.org" may use IPv6

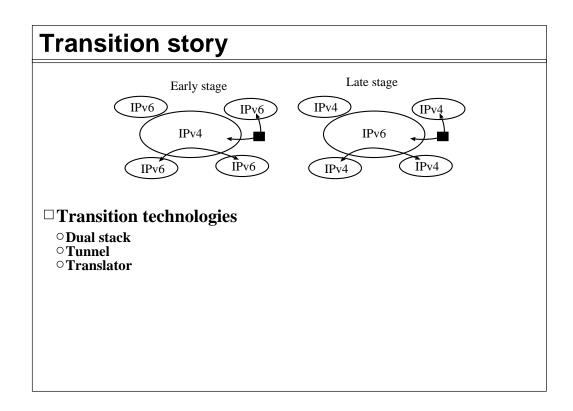
□ Users can specify an IPv6 address to applications, of course °(e.g) ftp 3ffe:501:8:1234:260:97ff:fe40:efab °(e.g) telnet ::1

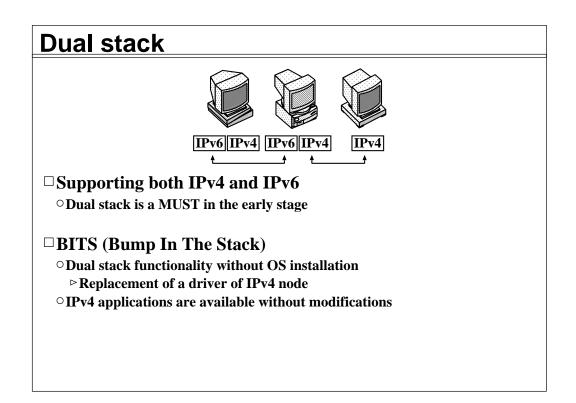
□ But '':'' is unfriendly to some application syntax ° URL > http://[3ffe:501:8:1234:260:97ff:fe40:efab]/

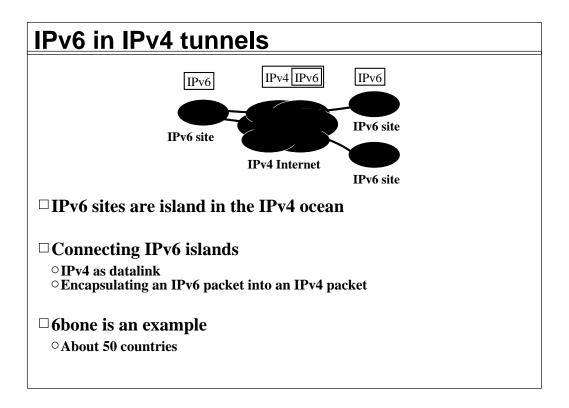
Domain Name System

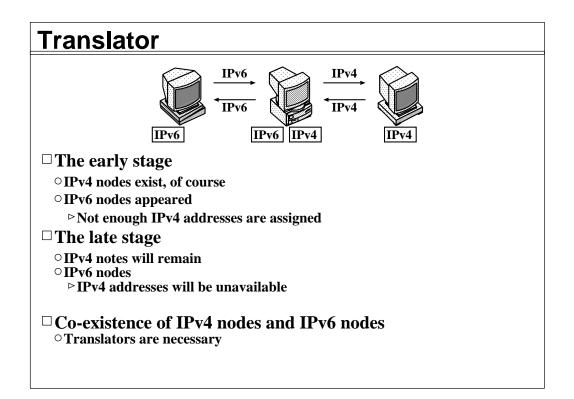


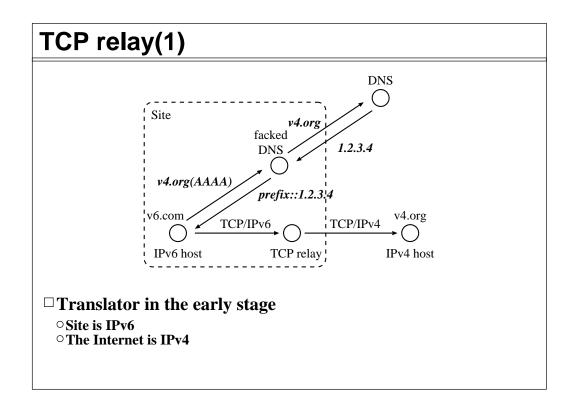












TCP relay(2)

□ **Precondition**

• IPv6 site has also small number of IPv4 addresses

• Both fake DNS and TCP relay are dual stack

▷A special prefix is installed on them statically

□ Connection from v6.com to v4.org

°v6.com asks faked DNS to resolve AAAA record of v4.org

°Faked DNS asks DNS to resolve A record of v4.org

ODNS returns 1.2.3.4

 $^{\rm O}$ Faked DNS embeds 1.2.3.4 into special prefix and tells v6.com it

 $^{\rm O}$ v6.com tries to make a TCP/IPv6 connection to prefix::1.2.3.4

°TCP relay catches the connection

°TCP relay also make a TCP/IPv4 connection to 1.2.3.4

What's IPv6 Address Architecture Plug & play Domain Name System Transition

=> The current status of IPv6

Specification and implementation

Basic specifications have been done

Some advanced specifications are being discussed

Many UNIX systems are IPv6-ready

BSD, Linux, Solaris 8, DEC(Compaq), IBM...

Windows 2000 + IE

http://msdn.microsoft.com/downloads/sdks/platform/tpipv6.asp

MacOS X will be ready in this year

Based on KAME

Many open source servers/clients are IPv6 ready

Sendmail, BIND, Apache, ...

Routers

Cisco, Hitachi, Fujitsu, NEC, ...

Deployment status

□ Research network ◦ 6bone (spans across 50 countries)

□ Experimental service of commercial ISP °IIJ, NTT Com, vBNS, NEC, Fujitsu

□ IX (Internet Exchange) ○NTT IX, 6TAP, PAIX, NSPIXP6

KAME Project

□A single effort

08 core members from 7 Japanese companies

°Fujitsu, Hitachi, IIJ, NEC, Toshiba, YDC, Yokogawa

○April 1998 - March 2002

 $^{\rm O}$ The core members work for IPv6 three days a week

□ Reference code

○ IPv6, IPsec, and advanced networking
 ○ Provided "AS IS" like BSD
 ▷ Free and no warranty, commercial use is OK

□Adopted

BSD/OS 4.2, FreeBSD 4.1, NetBSD 1.5 (2000 Autumn),
OpenBSD 2.7, IIJ SEIL T1, Hitachi GR2000, Fujitsu NetVehicle

TAHI Project

• Conformance test, inter-operability test

Information

ohttp://www.ipv6.org/ ohttp://www.6bone.net/ ohttp://www.v6.wide.ad.jp/ ohttp://www.kame.net/ ohttp://www.v6forum.com/ ohttp://www.freenet6.net/ ohttp://www.6tap.net/

$\Box ML$

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Ongtrans@sunroof.eng.sun.com
Oipng@sunroof.eng.sun.com
Ousers@ipv6.org

Global IPv6 Summit

Dec 18 - 19, 2000
Osaka, Japan
http://www.jp.ipv6forum.com

Program

- $^{\rm O}$ Keynote speech from Jun Murai and Steve Deering
- ○Business report from Japan
- ^OBusiness report around the world
- °Status report from CN, KR, SG, MY
- **OAddress allocation/assignment session**
- Panel on how IPv6 change business
- Panel on how to transit to IPv6